

The condition for $f(n) = O(g(n))$ is
 $f(n) < C g(n) \quad \forall n > n_0$

which can be restated

$$\frac{f(n)}{C g(n)} < 1 \quad \forall n > n_0 \quad (1)$$

(a) $f(n) = O(g(n))$ if

$$f(n) \leq C g(n) \quad \forall n > n_0$$

for arbitrary C and n_0

$f(n) = \Omega(g(n))$ if

$$f(n) \geq C g(n) \quad \forall n > n_0$$

for arbitrary C and n_0 .

Note: the C and n_0 need not be the same
clearly

$$f(n) \leq C g(n) \quad \forall n > n_0$$

implies

$$g(n) \geq \frac{1}{C} f(n) \quad \forall n > n_0$$

So $g(n) = \Omega(f(n))$ if $f(n) = O(g(n))$

$$(b) \quad f(n) = \sum_{j=0}^k a_j n^j \quad g(n) = n^l$$

$$\therefore \frac{f(n)}{g(n)} = \sum_{j=0}^k a_j n^{j-l}$$

If $l > k$ then all n^{j-l} for $j \leq k$ will vanish as $n \rightarrow \infty$, so

$$\frac{f(n)}{g(n)} \rightarrow 0 \text{ as } n \rightarrow \infty$$

$\therefore f(n) = o(n^l)$ by Eg. 1

$$(c) \quad \text{If } \frac{a(n)}{c g(n)} < 1 \quad n > n_0$$

$$\frac{b(n)}{c' h(n)} < 1 \quad n > n'_0$$

$$\text{then } \frac{a(n)}{c g(n)} \frac{b(n)}{c' h(n)} < 1 \quad n > \max(n_0, n'_0)$$

$$\therefore \frac{a(n)b(n)}{c c' g(n) h(n)} < 1 \quad n > \max(n_0, n'_0)$$

$$\therefore a(n)b(n) = o(g(n) h(n)) \text{ if } a(n) = o(g(n)) \text{ and } b(n) = o(h(n))$$

The test are just more of the same