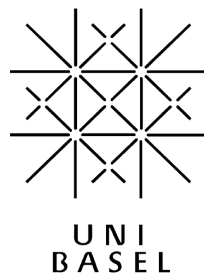
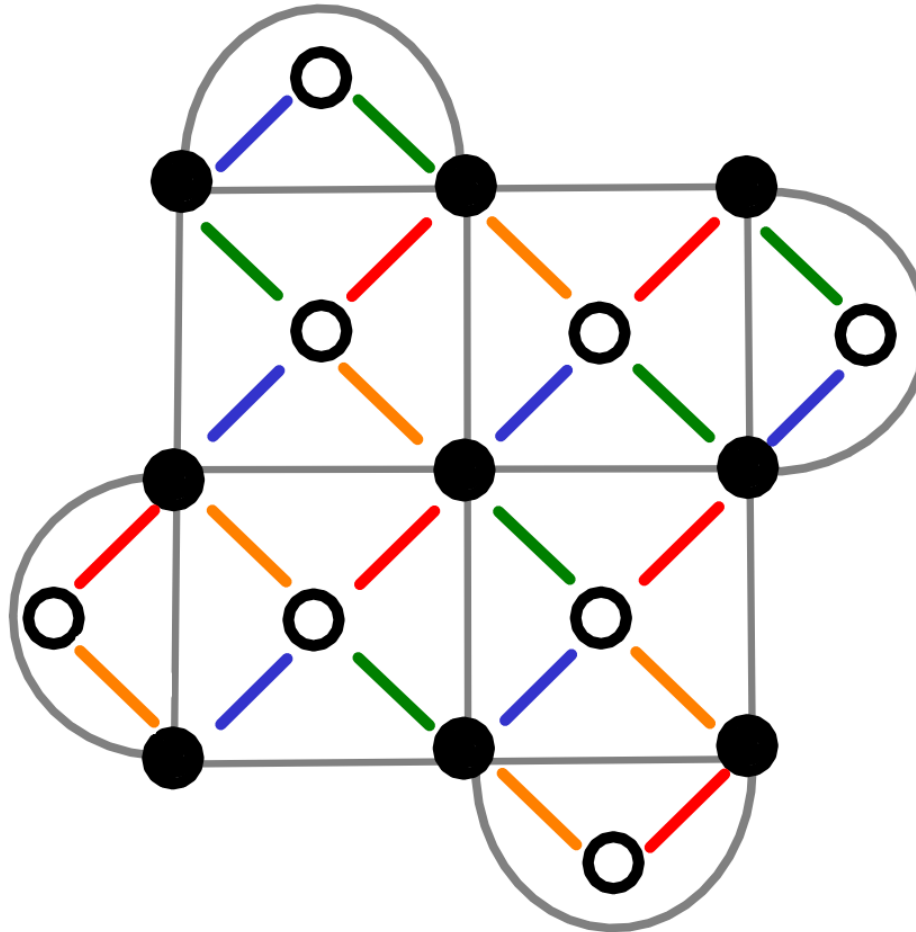


Introduction to the Surface Code

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Science and
Technology*

National Centre of Competence in Research

Towards a better quantum code

- How does the repetition code protect against bit flip noise (σ_x) ?

- An isolated σ^x creates a pair of *defects*

0 0 0 ≠ 1 ≠ 0 0 0 0 0 0 0

- Further σ^x s can move the defects

0 0 0 ≠ 1 1 ≠ 0 0 0 0 0 0

- Or create new pairs of defects

0 0 0 ≠ 1 1 ≠ 0 ≠ 1 ≠ 0 0 0 0

- Or annihilate pairs of defects

0 0 0 ≠ 1 1 1 1 ≠ 0 0 0 0

- A distance of $>d/2$ is needed for a logical error

0 0 0 ≠ 1 1 1 1 1 1 ≠ 0 0

- Then decoding will complete the job, pulling the defects off the ends

≠ 1 1 1 1 1 1 1 1 1 1 ≠

- The code is like a ‘universe’ in which the defects are its particles

- Bit flips create and manipulate these particles, but only large scale effects can cause a logical error

Towards a better quantum code

- Why doesn't the repetition code protect against phase flip noise (σ_z) ?
- Measurement is too easy, even when the information is encoded

0 0 0 0 0 0 0 0 0 0 0 0

1 1 1 1 1 1 1 1 1 1 1 1

- Once errors are removed, a quick peek at any qubit reveals the stored information
- If it is easy for us to see, it is easy for the environment to dephase
- Consider measuring in the X basis instead

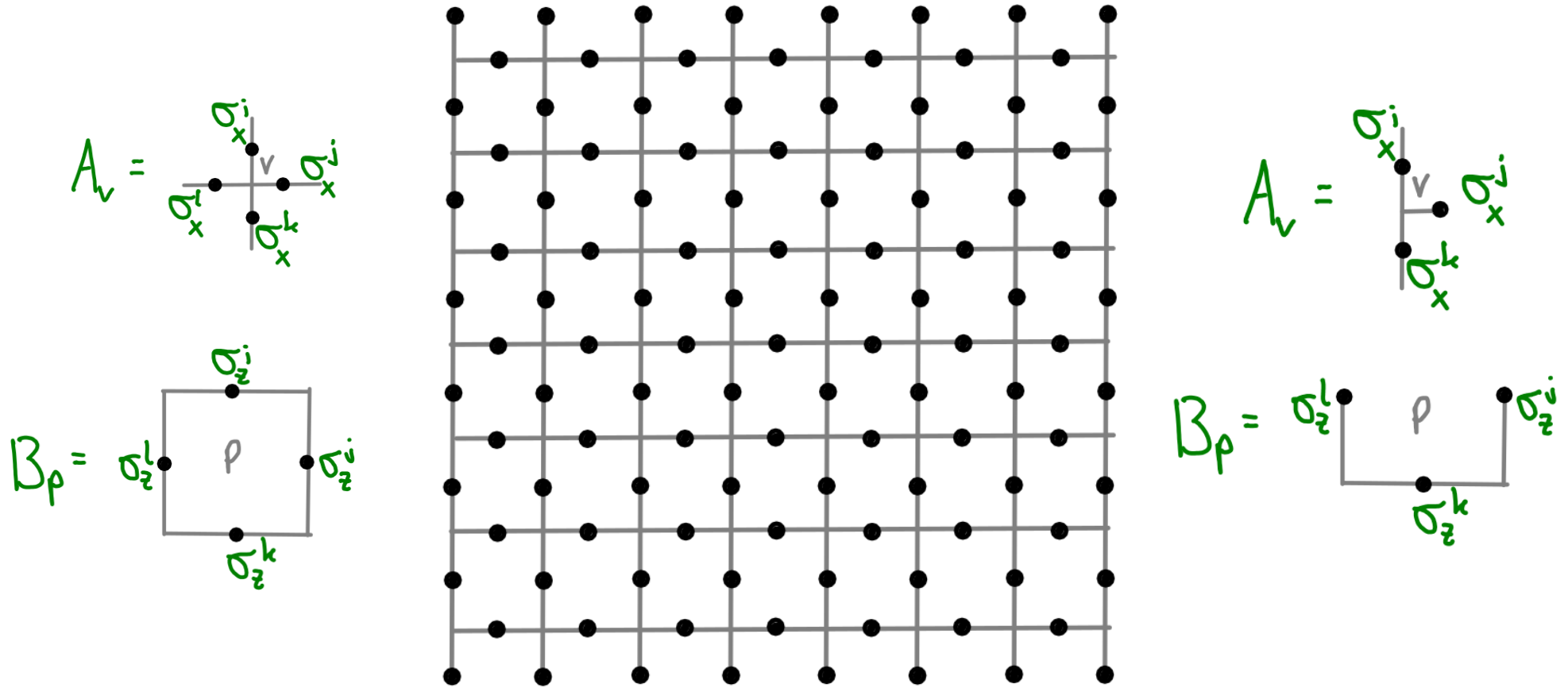
$$|\pm\rangle = \frac{1}{\sqrt{2}}(|0\rangle \pm |1\rangle) \rightarrow \frac{1}{\sqrt{2}}(|000\rangle \pm |111\rangle)$$

- Requires multi qubit process for the encoded states

Towards a better quantum code

- This code treats the X and Z basis of the qubit very differently
- We need a code that treats them the same
 - ✓ σ_x creates particle-like defects that can be detected
 - ✓ Large scale effects are needed for a logical bit flip
 - ✓ Multiqubit measurement needed to distinguish encoded $|+\rangle, |-\rangle$
 - ✓ σ_z creates particle-like defects that can be detected
 - ✓ Large scale effects are needed for a logical phase flip
 - ✓ Multiqubit measurement needed to distinguish encoded $|0\rangle, |1\rangle$
- Other methods of generalizing the repetition code also exist, like the Shor code
- But these don't create new universes, and are therefore boring

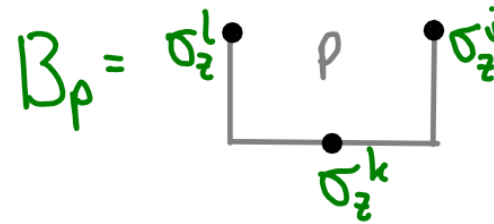
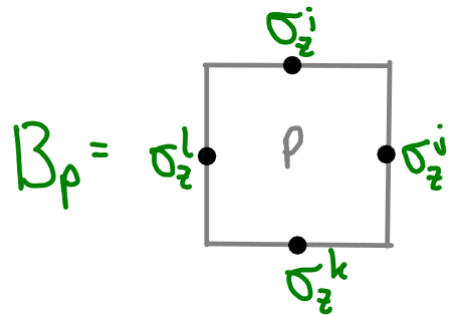
The Surface Code



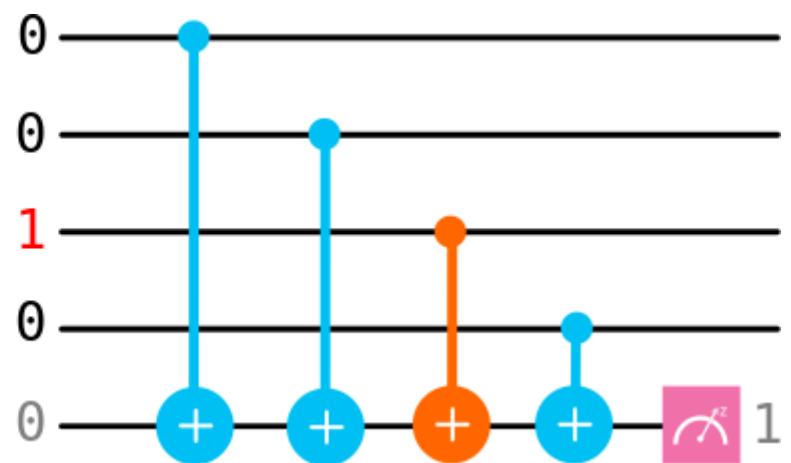
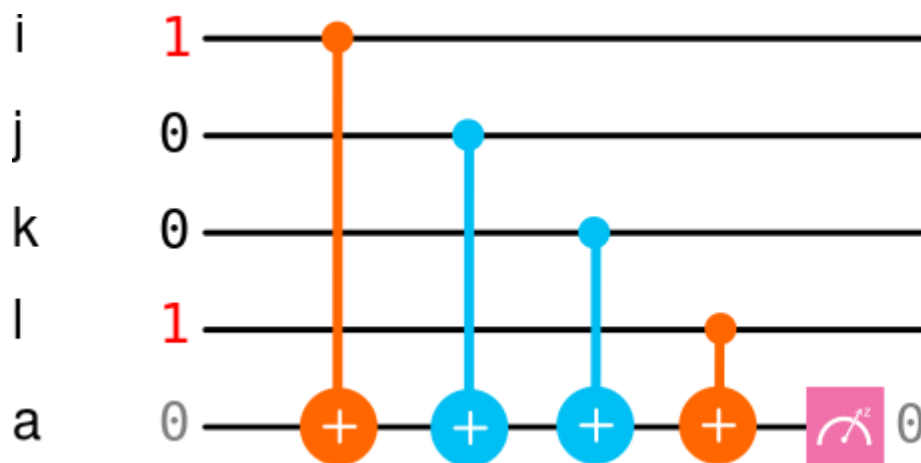
- To do this, we put our grid on a 2D lattice
- The $\sigma_z^j \sigma_z^{j+1}$ observables between neighbouring qubits become ones for qubits around plaquettes
- Similar observables for σ_x are defined on vertices

The plaquette operators

- Let's focus on the plaquette operators



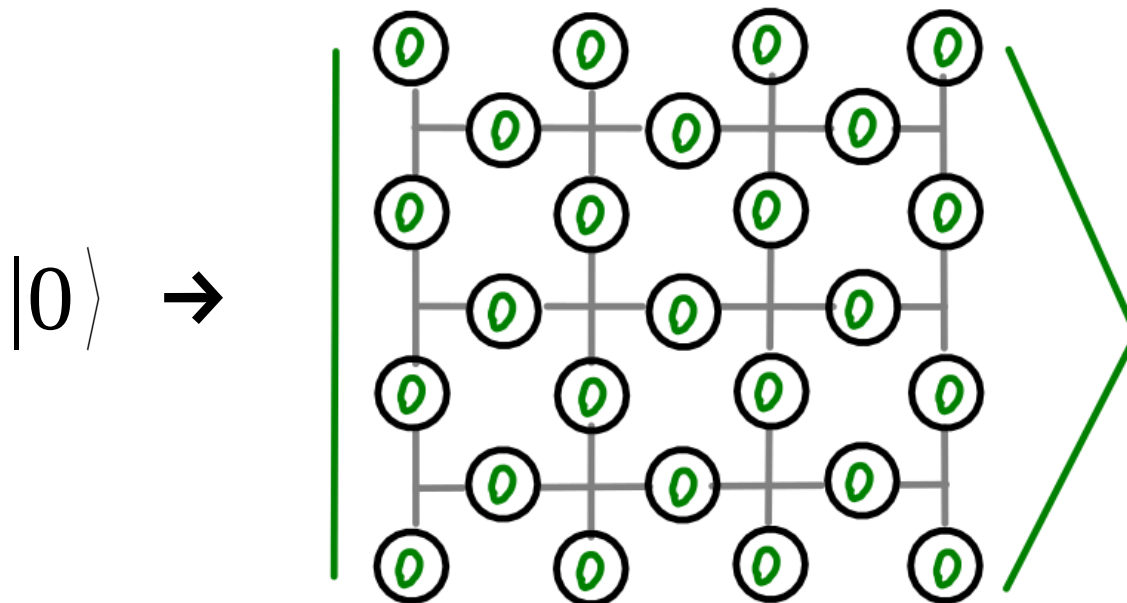
- Generalization of the measurement in the repetition code
- Can be similarly implemented with the controlled-NOT



- They tell us whether there is an odd or even # 1s around the plaquette

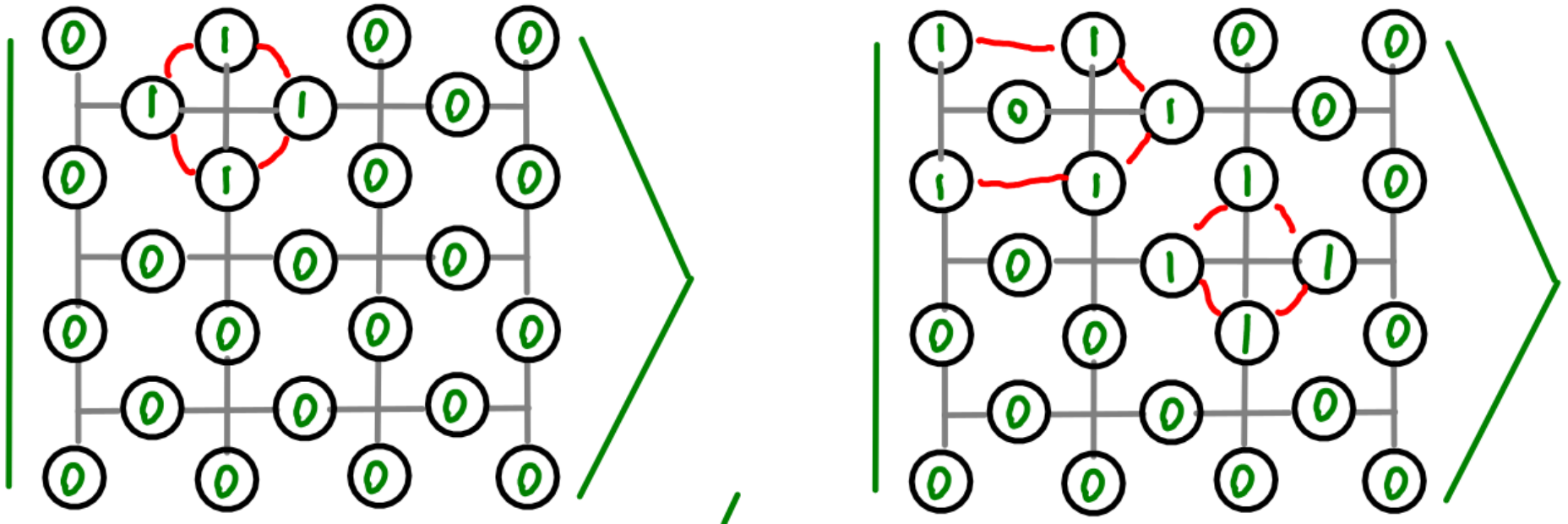
The plaquette operators

- How do we store a bit in this code?
- Valid encoded states are those for which the measurements don't detect an error
- We associate this with outcome 0, so all plaquettes need an even # 1s
- Let's again encode 0 with the 'all qubits are 0' state



The plaquette operators

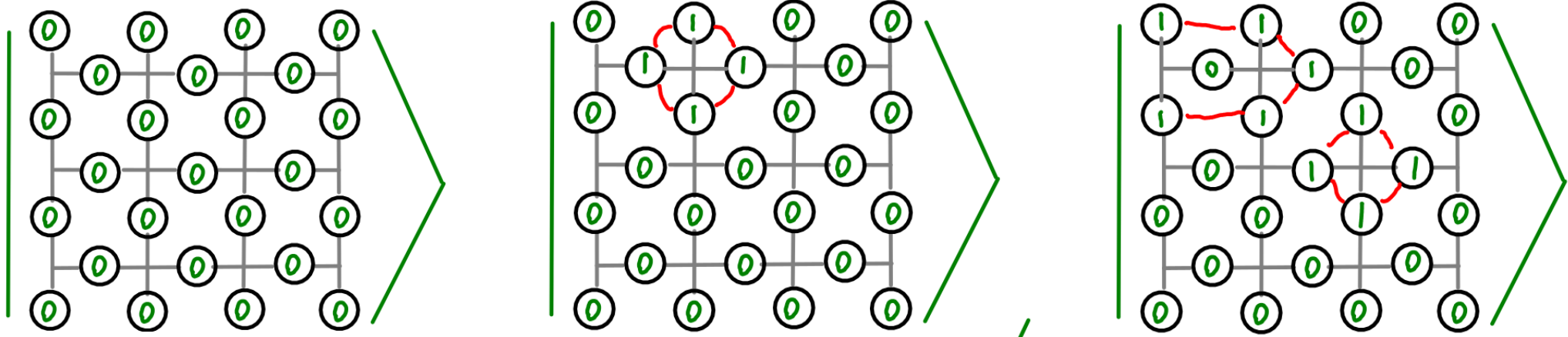
- There are other 'nearby' states that have the same results for plaquette measurements



- They can't be our encoded 1, because they differ by only a few bit flips
- So let's treat them as other possible ways to encode 0

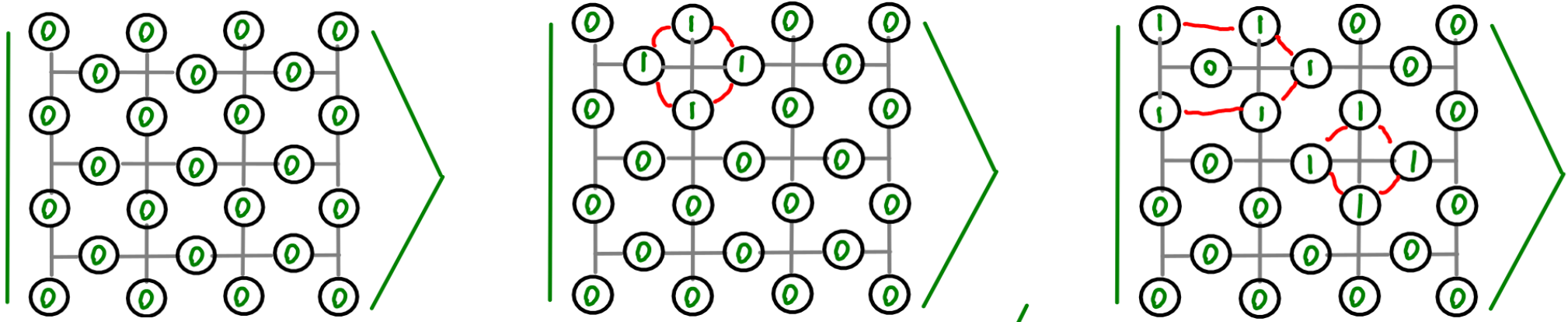
Encoding 0 and 1

- Given any possible encoding for 0:
 - 1) Pick a vertex
 - 2) Apply a bit flip on all qubits around the vertex
- Now you have another possible encoding for 1
- This generates an exponentially large family



Encoding 0 and 1

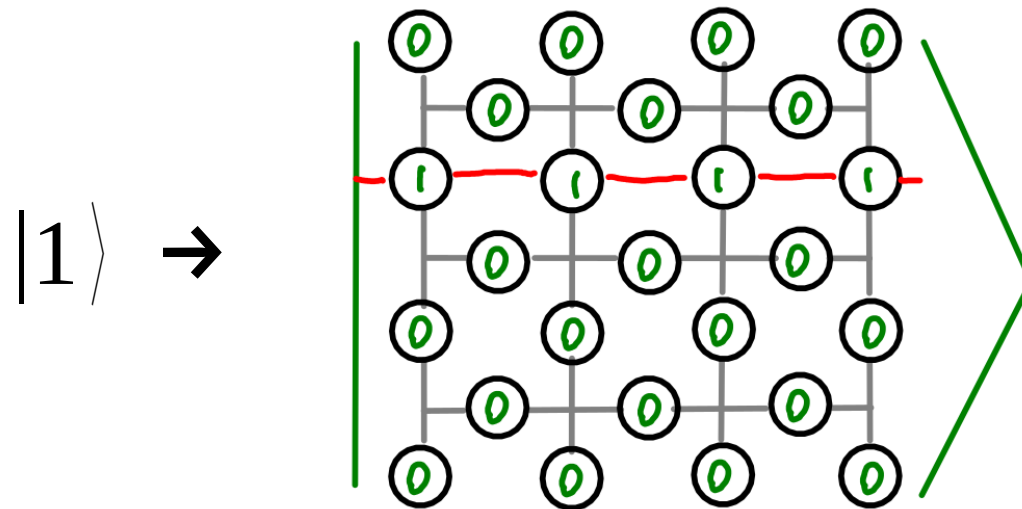
- The states in this family can be very different
- But there is one feature shared:
A line from top to bottom will always have an even number of 1s
- This is how we can measure our encoded 0 state



- And it gives us a hint on how to encode a 1

Encoding 0 and 1

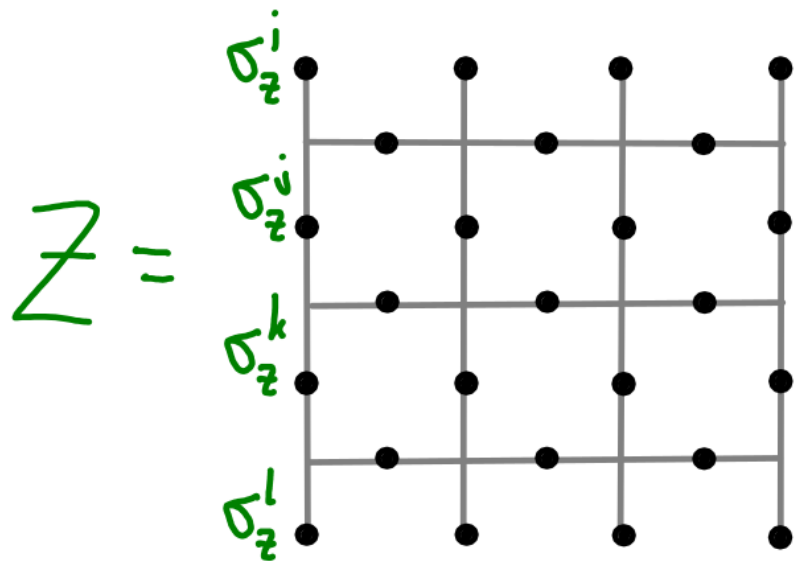
- As our basic encoded 1, we can use a bunch of 0s with a line of 1s across



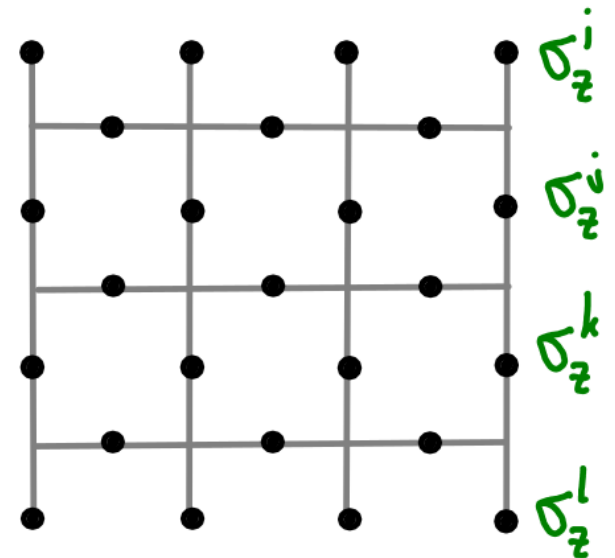
- This also spawns an exponentially large family
- For each state in that family, the number of 1s on a line from top to bottom is odd
- Measuring our encoded bit has become hard (which is good!)

X and Z for encoded qubit

- Measuring 0 and 1 corresponds to measuring an observable Z for the encoded qubit
- The observable that detects what we need is



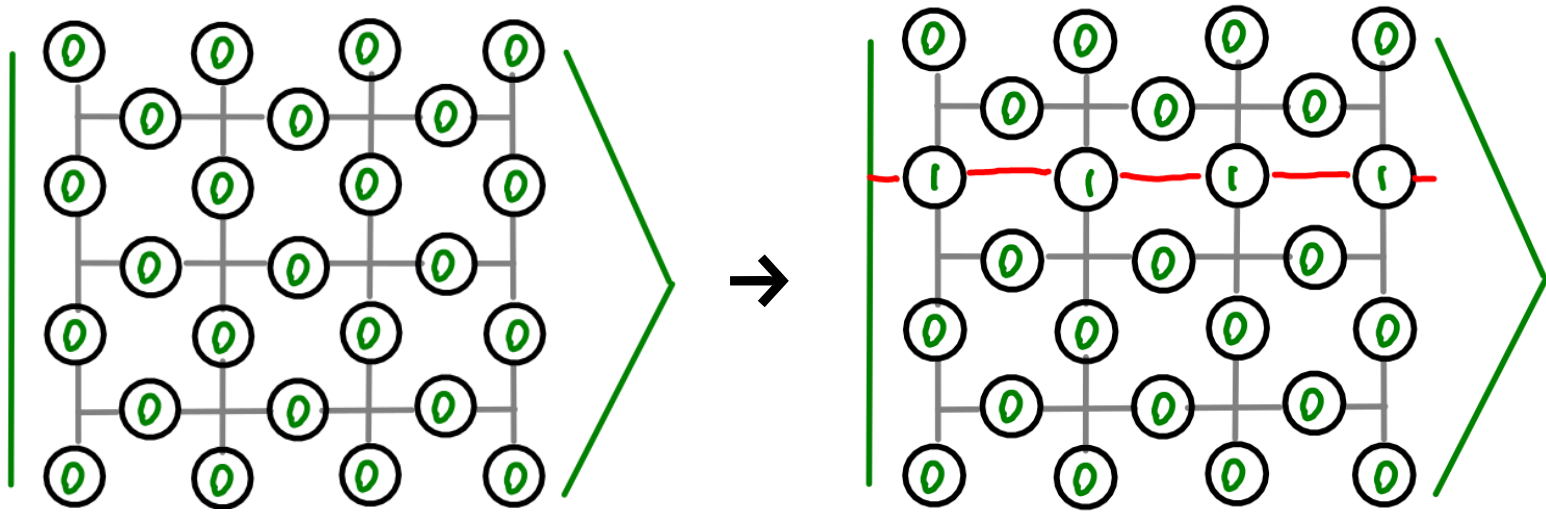
$Z' =$



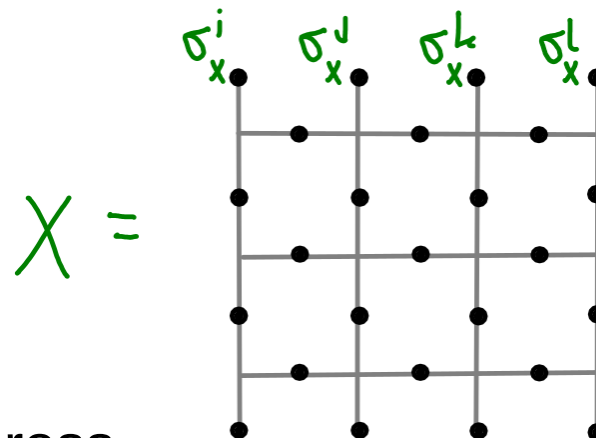
- Or the same on any line from top to bottom
- If we use the edges, we can think of them as large and unenforced plaquettes

X and Z for encoded qubit

- If we want to do a bit flip on the encoded bit, clearly we need a line of flips from left to right



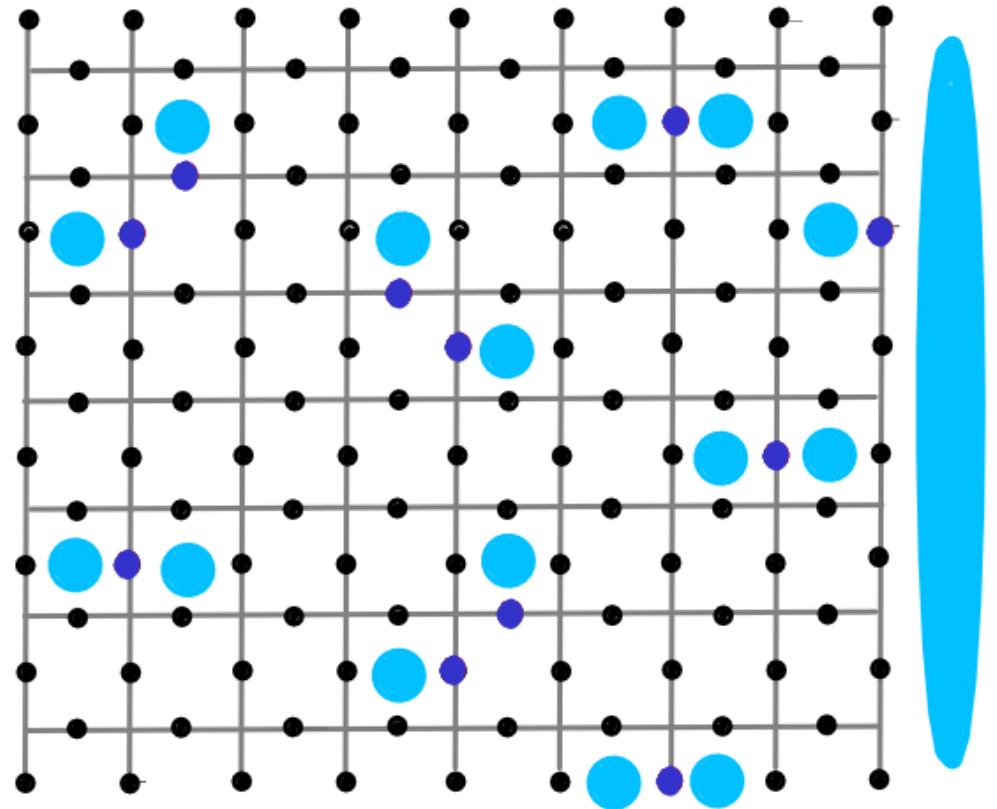
- So the X operation for an encoded qubit is



- Or the same on any other line across

Effects of errors

- What happens when a σ_x is applied?
- Changes measurement outcome for neighbouring plaquettes
 - An isolated σ_x creates a pair of defects
 - Further σ_x s can move the defects
 - Or create new pairs of defects
 - Or annihilate pairs of defects
 - A distance of $>d/2$ is needed for a logical error, where d is the width



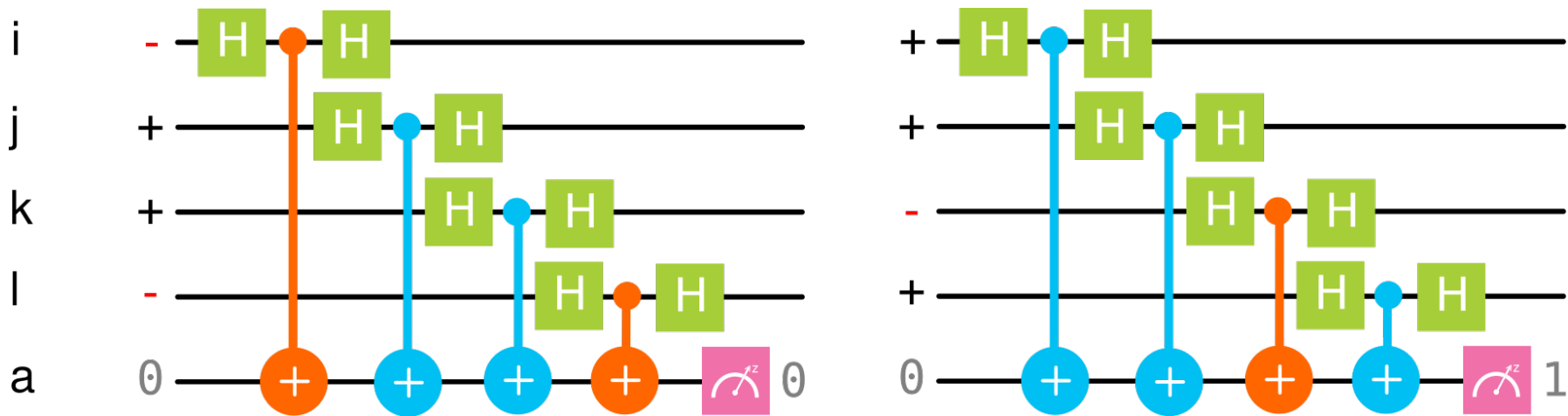
- With the plaquette operators we can encode and protect a bit using Z basis states

Vertex operators

- Now forget the plaquettes, and focus on the vertices

$$A_v = \begin{array}{c} \sigma_x^i \\ \bullet \\ \sigma_x^j \\ \bullet \\ \sigma_x^k \\ \bullet \\ \sigma_x^l \end{array} \quad A_v = \begin{array}{c} \sigma_x^i \\ \bullet \\ \sigma_x^j \\ \bullet \\ \sigma_x^k \\ \bullet \\ \sigma_x^l \end{array}$$

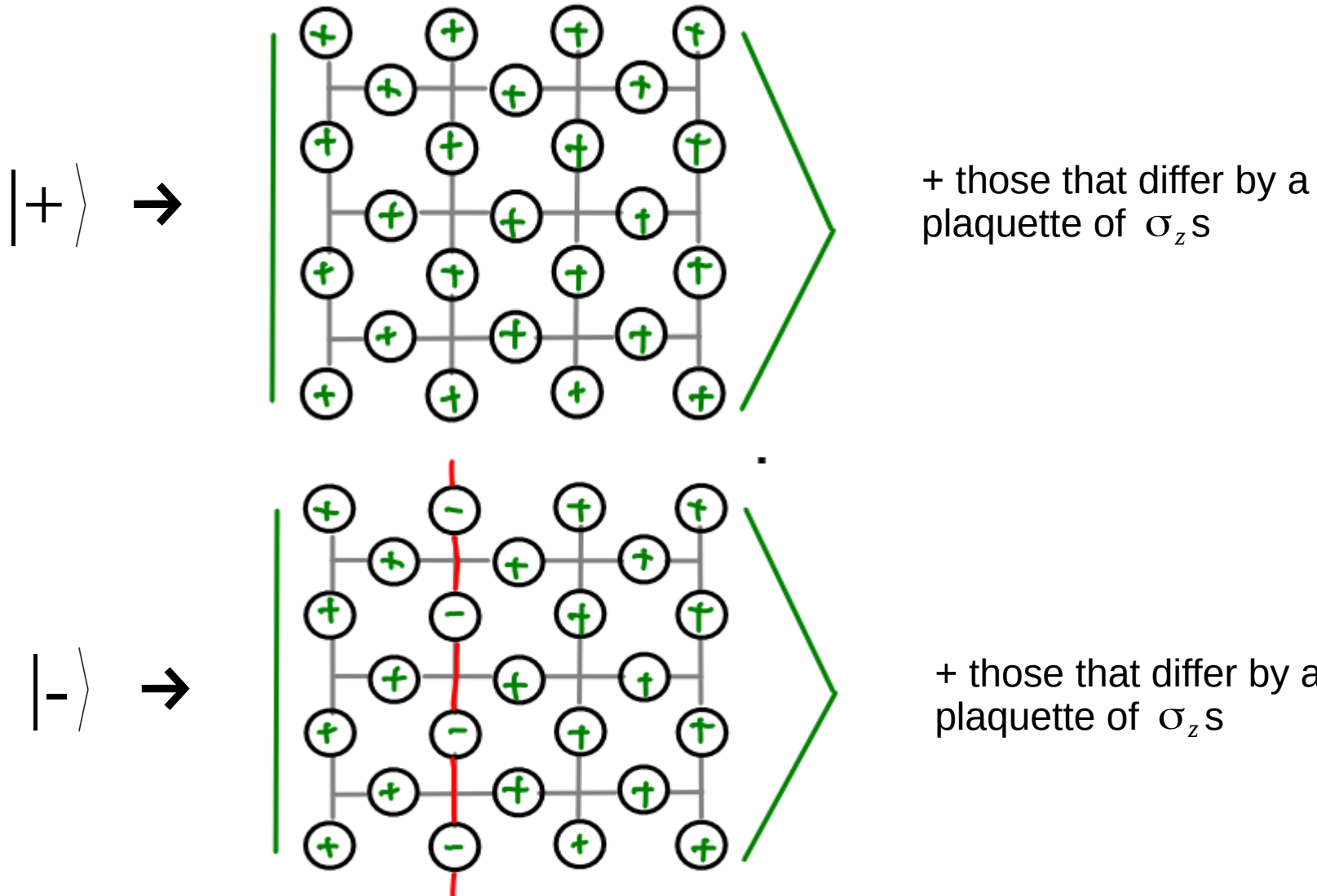
- These can also be measured with controlled ops and an ancilla



- Looks at $|+\rangle$ and σ^x states, and tell us whether there is an even number of $|-\rangle$ s

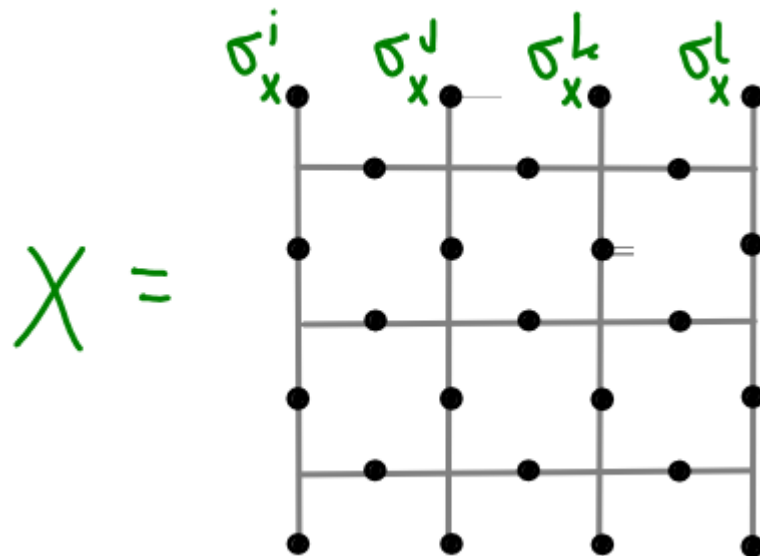
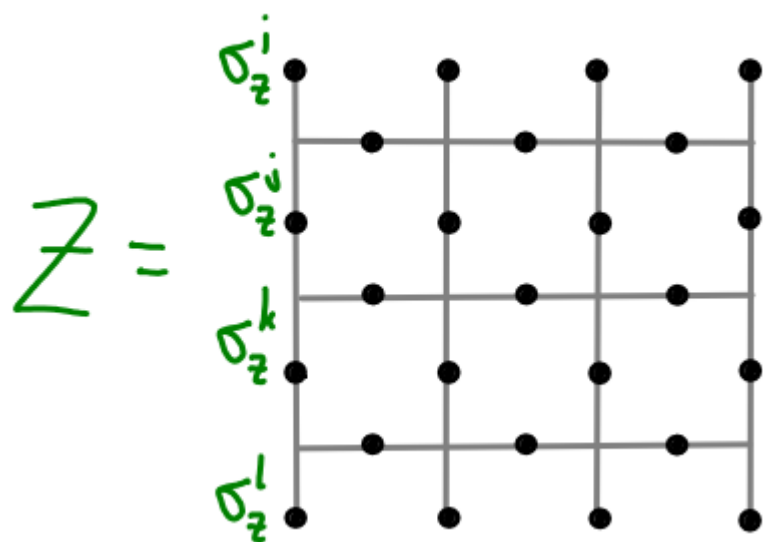
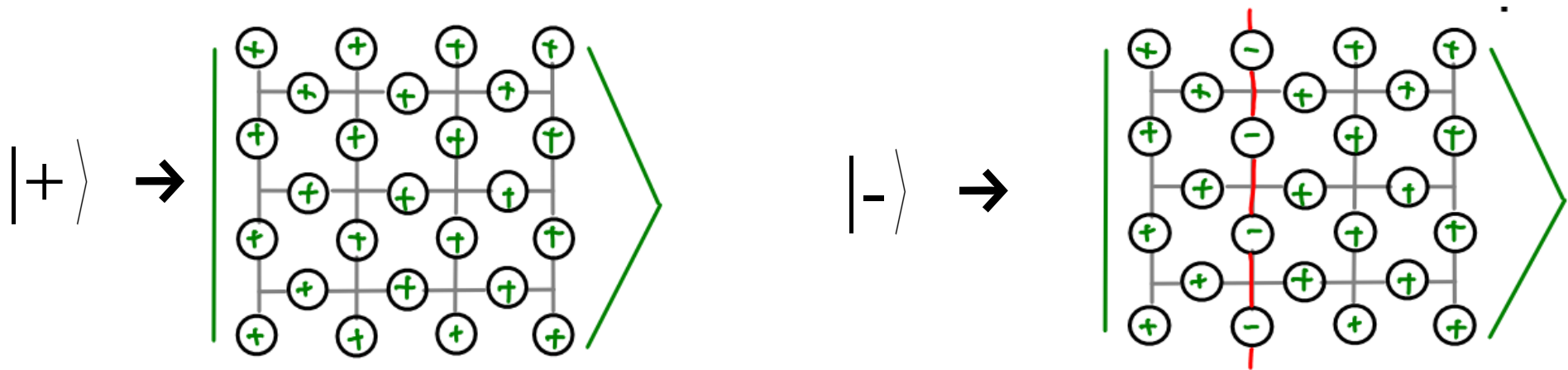
Encoding + and -

- We can also store a bit using only the vertex ops
- Let's associate this with the x basis instead



Encoding + and -

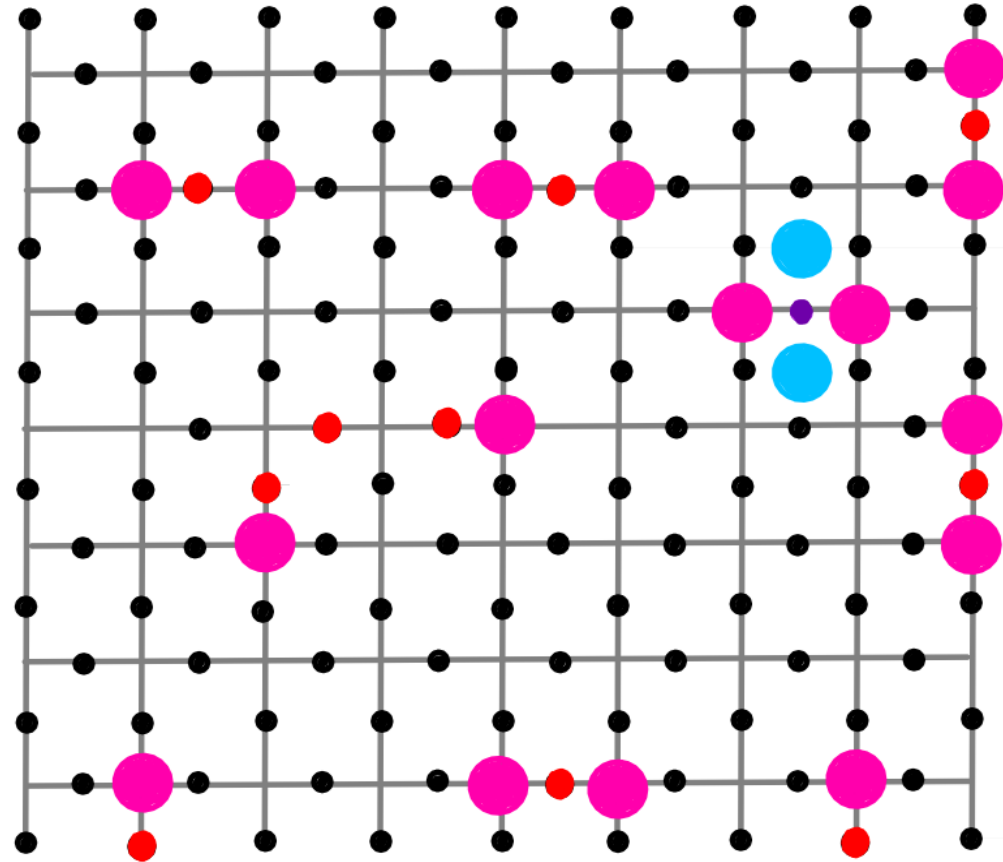
- This leads to exactly the same logical operators as before



Effects of errors

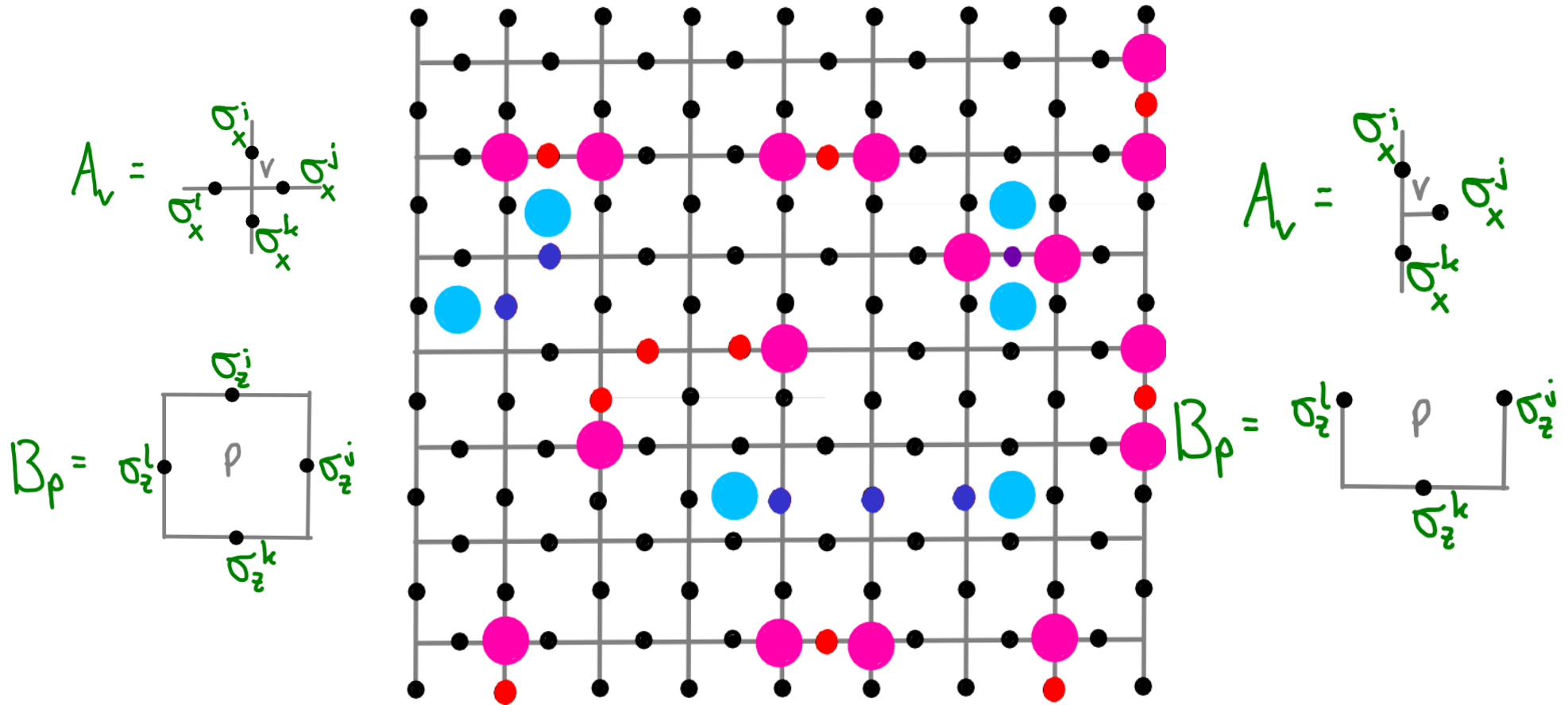
- What happens when a σ_z is applied?
- Changes measurement outcome for neighbouring plaquettes

- An isolated σ_z creates a pair of *defects*
- Further σ_z s can move the defects
- Or create new pairs of defects
- Or annihilate pairs of defects
- A distance of $>d/2$ is needed for a logical error, where d is the height



- With the vertex operators we can encode and protect a bit using X basis states

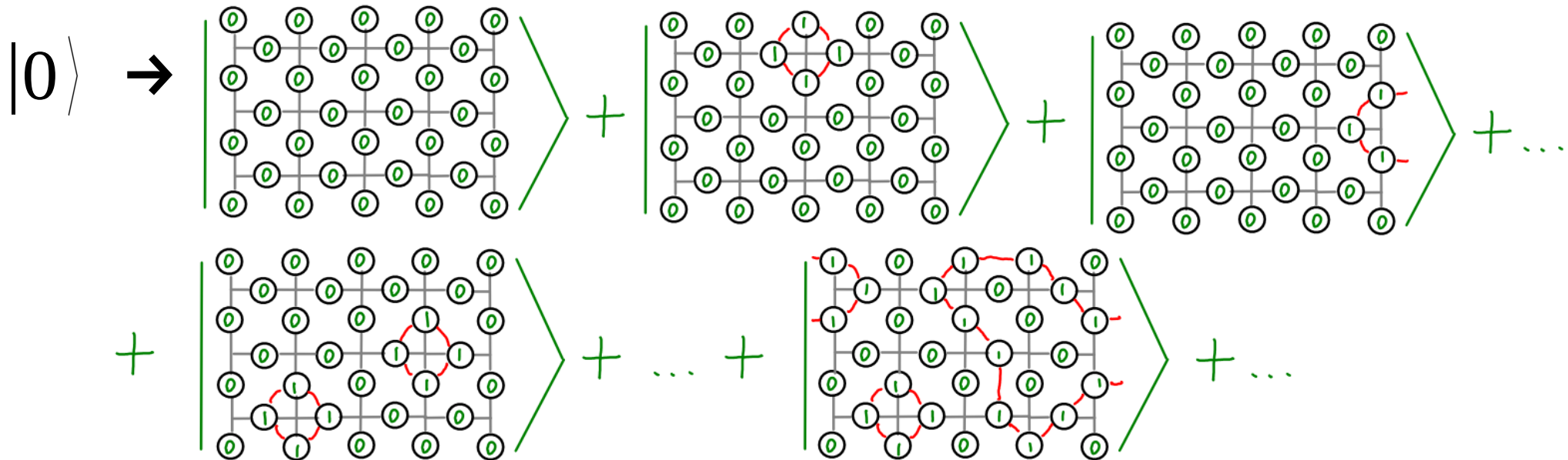
Vertex and plaquette operators together



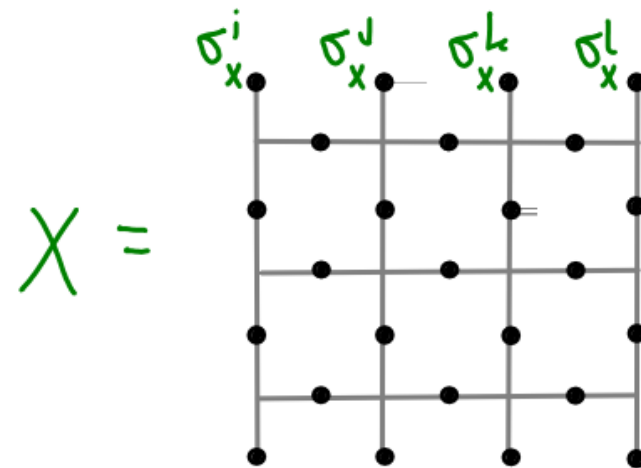
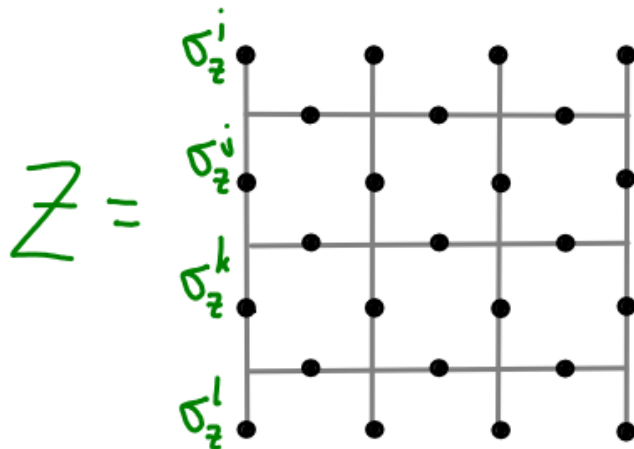
- The vertex and plaquette operators commute
- We can measure these observable simultaneously
- Detect and correct σ_x and σ_z errors simultaneously

Vertex and plaquette operators together

- Encoded states now unique: superposition of all previous solutions
- Highly entangled states

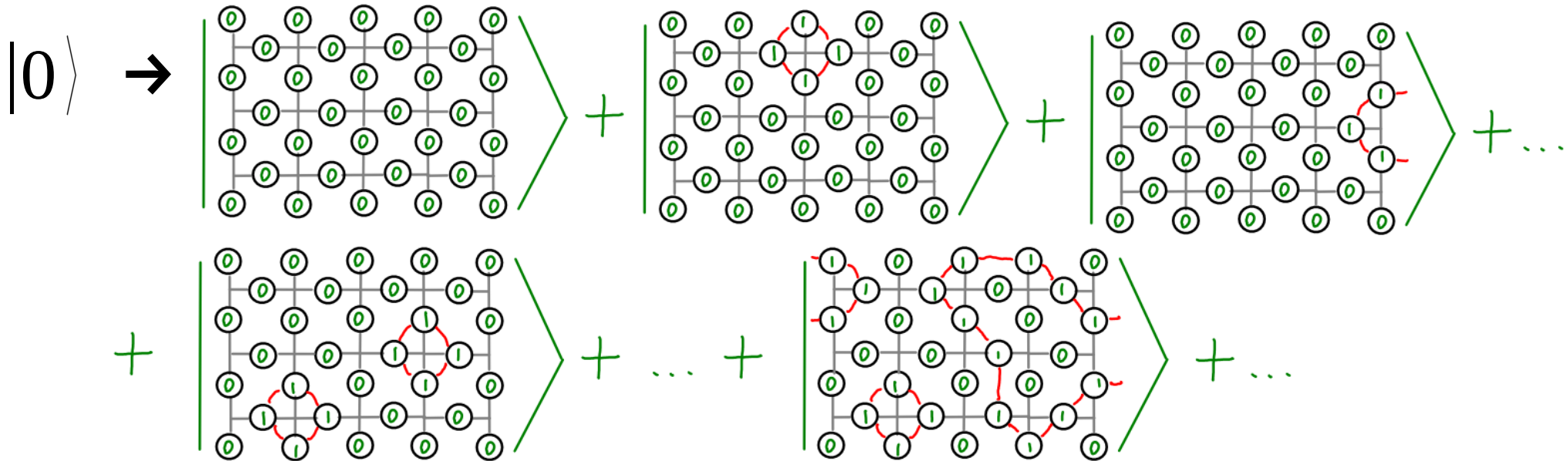


- Pauli X and Z for encoded qubits exactly as they were for plaquettes and vertices alone

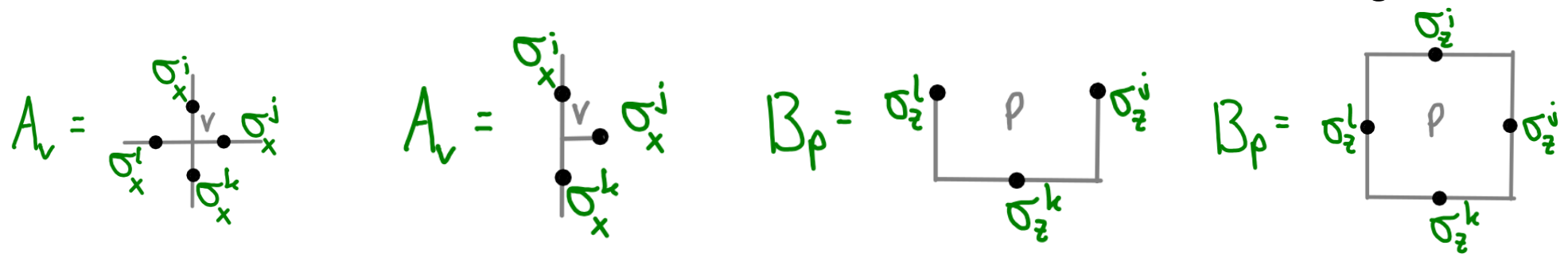


Vertex and plaquette operators together

- But aren't many-body entangled states hard to make?



- They are the mutual eigenstates of the observables we measure
- If we can measure them, we can create and maintain the entanglement



Vertex and plaquette operators together

- We are not just protected against σ_x and σ_z noise, but all local noise
- Any noise operator can be expressed in terms of Paulis

$$M |\psi\rangle = a \sigma_0 + b \sigma_x + c \sigma_y + d \sigma_z$$

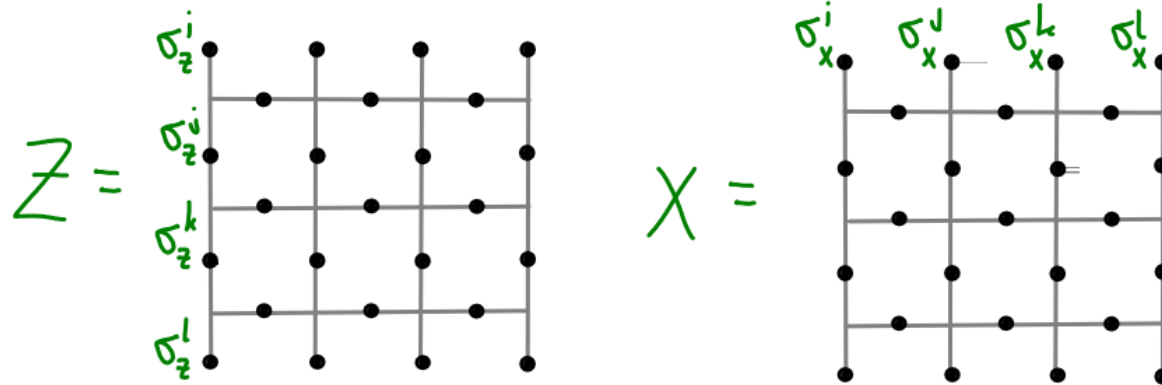
- And so creates a superposition of different measurement outcomes for the plaquettes and vertices

$$M |\psi\rangle = a \text{ [grid with purple dot] } + b \text{ [grid with blue dots] } + c \text{ [grid with pink and blue dots] } + d \text{ [grid with pink dots] }$$

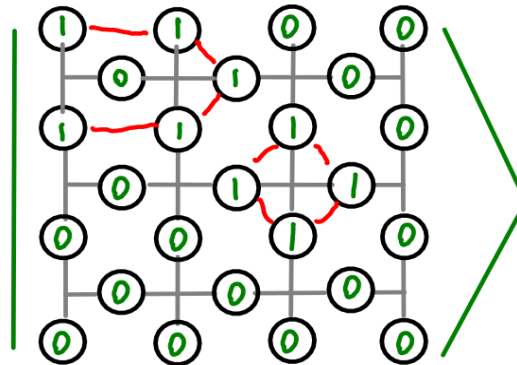
- Measurement collapses this superposition, reducing noise to a simple Pauli
- So any noise can be detected and corrected

Final Readout

- The logical operators are many body operations



- How do we read out stored information without error?
- When you decide on a basis, you stop caring about one kind of error
- We can just measure in a product basis



- Logical Z and plaquette info can be constructed from the result
- Imperfect measurement can be corrected like a bit flip

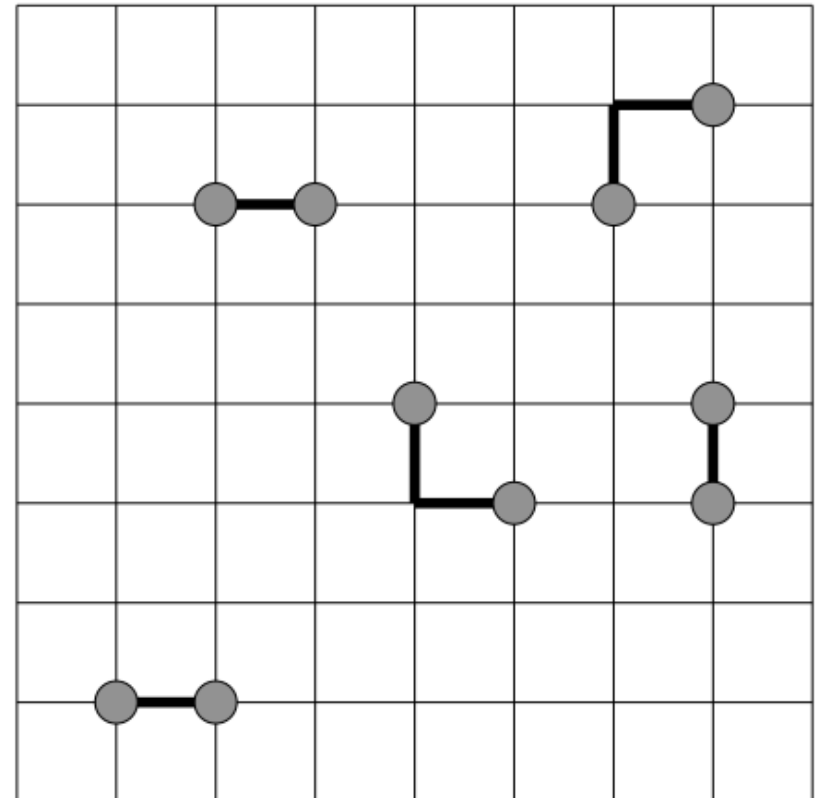
Imperfect measurement

- What about imperfect measurements throughout?
- Consider a measurement of a single qubit that lies with prob. P , but doesn't disturb the measured qubit (beyond projection)
- How do we extract information correctly? Repetition!
- Lies create pairs of defects in the time direction



Imperfect measurement

- Combine this with the repetition code or surface code:
 - Defects = changes in ancilla measurement result
 - Bit flips create space-like separated defect pairs
 - Lies create time-like separate defect pairs
 - Combinations create combinations
- Noisy measurements just increase the space of the 'universe' by 1 dimension



The surface code is a good quantum code

- X and Z basis are treated the same
 - ✓ σ_x creates particle-like defects that can be detected
 - ✓ Large scale effects are needed for a logical bit flip
 - ✓ Multiqubit measurement needed to distinguish encoded $|+\rangle, |-\rangle$
 - ✓ σ_z creates particle-like defects that can be detected
 - ✓ Large scale effects are needed for a logical phase flip
 - ✓ Multiqubit measurement needed to distinguish encoded $|0\rangle, |1\rangle$
- Other good quantum codes also exist
 - Topological codes: Color code, quantum double modes, ...
 - Concatenated codes: Shor code, ...